

# CoS for MPLS VPNs with IP Quality of Service

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### CoS vs. QoS

Class of Service (CoS) is the service provided by a network to the customer traffic.

In IP based networks, this service is achieved by using IP Quality of Service (QoS) mechanisms and features.

# More Bandwidth? Is that always the Answer?



#### **PROS**

- Increases capacity
- Resolves immedia

#### CONS

- Short-term solution
- Expensive \$\$\$
- Will not guarantee applications with low latency tolerance such as VoIP and video conferencing
- All applications receive same service, no protection for mission-critical applications
- Emerging applications could jeopardize business critical traffic



## How Can QoS Be Applied?

- Best effort—no QoS is applied to packets (default behavior)
- Integrated Services model—applications signal to the network that they require special QoS
- Differentiated Services model—the network recognizes classes that require special QoS



# Integrated Services Model (Intserv)

# Benefits and Drawbacks of the IntServ Model



#### + RSVP benefits:

- Explicit resource admission control (end-to-end)
- Per-request policy admission control (authorization object, policy object)
- Signaling of dynamic port numbers (for example, H.323)

#### -RSVP drawbacks:

- Continuous signaling due to stateless architecture
- Not scalable
- Needs modified (RSVP supporting) applications



# Differentiated Services Model (DiffServ)



### Differentiated Services Model

- TheDifferentiated Services model describes services associated with traffic classes.
- Complex traffic classification and conditioning are performed at network edge, resulting in a per-packet Differentiated Services Code Point (DSCP).
- No per-flow/per-application state exists in the core.
- The core performs only simple "per-hop behaviors" on traffic aggregates.
- The goal is scalability.

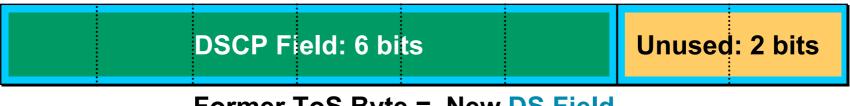


# Why Is Provisioning Important?

- QoS does not create bandwidth!
- QoS manages bandwidth usage among multiple classes.
- QoS gives better service to a well-provisioned class with respect to another class.



### Packet Header Terminology



Former ToS Byte = New <u>DS Field</u>

- DSCP: a specific value of the DSCP portion of the DS field. The DSCP is used to select a PHB (Per-Hop Behavior; forwarding and queuing method)
- DS field: the IPv4 header ToS octet or the IPv6 traffic class octet when interpreted in conformance with the definition given in <u>RFC 2474</u>. The bits of the DSCP field encode the DSCP, while the remaining bits are currently unused.



### **DSCP Encoding**

- Three pools:
  - "xxxxxx0" Standard Action
  - "xxxx11" Experimental/Local Use
  - "xxxx01" EXP/LU (possible std action)
- Default DSCP: "000000"
- Default PHB: FIFO, tail-drop



### **DSCP** Usage

# DSCP selects per-hop behavior (PHB) throughout the network:

- Default PHB
- Class selector (IP Precedence) PHB
- Expedited forwarding PHB
- Assured forwarding PHB



# **Queuing Overview**

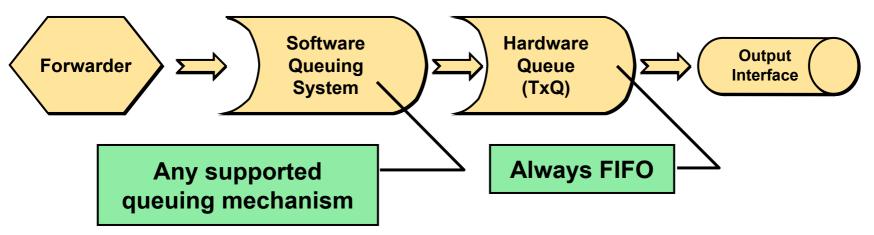


# **Queuing in Cisco IOS**

- Cisco routers running Cisco IOS have a number of different legacy queuing mechanisms
- Let's have a look first at the following
  - First In First Out (FIFO)
  - Weighted Fair Queuing (WFQ) with the different distributed versions
  - Priority Queuing (PQ)
  - Custom Queuing (CQ)
  - Modified Deficit Round Robin (MDRR)
  - IP RTP Prioritization
- These mechnisms are implemented as software queues



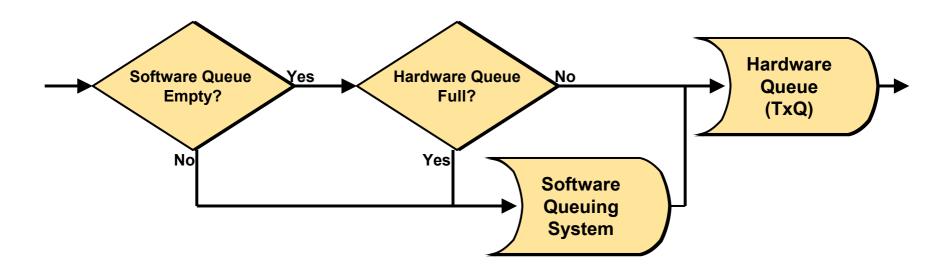
### **Output Interface Queue Structure**



- Each interface has its hardware and software queuing system.
- The hardware queuing system (transmit queue, or TxQ) always uses FIFO queuing.
- The software queuing system can be selected and configured depending on the platform and Cisco IOS version.



### **Bypassing the Software Queue**



- When a packet is being forwarded, the router will bypass the software queue if:
  - The software queue is empty, and
  - The hardware queue is not full



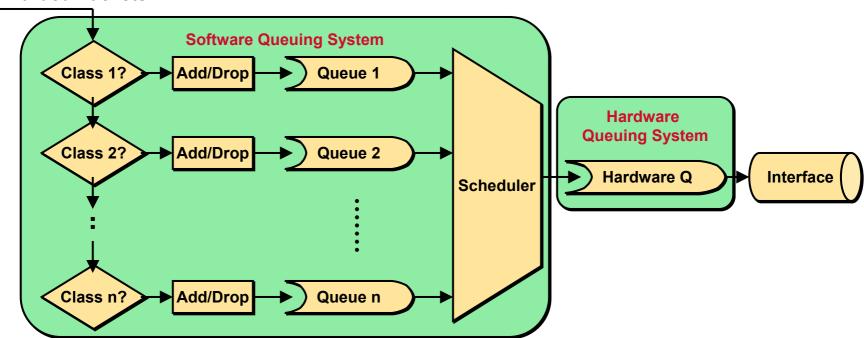
# Hardware Queue (TxQ) Size

- Routers determine the length of the hardware queue based on the configured bandwidth of the interface.
- Long TxQ may result in poor performance of the software queue.
- Short TxQ may result in a large number of interrupts which causes high CPU use and low link use.



## **Queuing Components**

#### **Forwarded Packets**



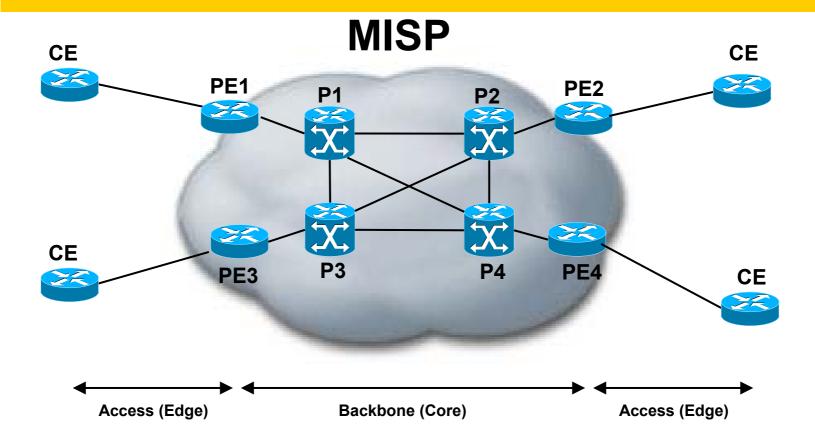
- Each queuing mechanism has three main components that define it:
  - Classification (selecting the class)
  - Insertion policy (determining whether a packet can be enqueued)
  - Service policy (scheduling packets to be put into the hardware queue)



# **CoS Implementation**

# Typical Service Provider Networks







### **QoS in Service Provider Networks**

- Service providers can extend their service offerings by introducing quality.
- Customers can get bandwidth guarantees (like CIR in Frame Relay).
- Customers can get delay guarantees (like CBR in ATM).
- QoS mechanisms have to be deployed where congestion is likely (usually at the network edge).
- The customer traffic is identified based on source or destination IP addresses or application type.



### Classes vs. DiffServ

# We offer Class 1, Class 2 and Class 3 services in three different combinations:

- Class 3 is best effort (default PHB)
  - DSCP 0 (default)
- Class 2 is Assured forwarding (AF PHB)
  - DSCP AF41 (34) for in-contract
  - DSCP AF43 (38) for out-of-contract
- Class 1 is Expedited forwarding (EF PHB)
  - **DSCP EF (46)**



### **Option (profile) A:**

- Class 3 gets 10% of available bandwidth
- Class 2 gets 50% of available bandwidth
- Class 1 gets 40% of available bandwidth with a low-delay guarantee



### **Option (profile) B:**

- Class 3 gets 60% of available bandwidth
- Class 1 gets 40% of available bandwidth with a low-delay guarantee



### **Option (profile) C:**

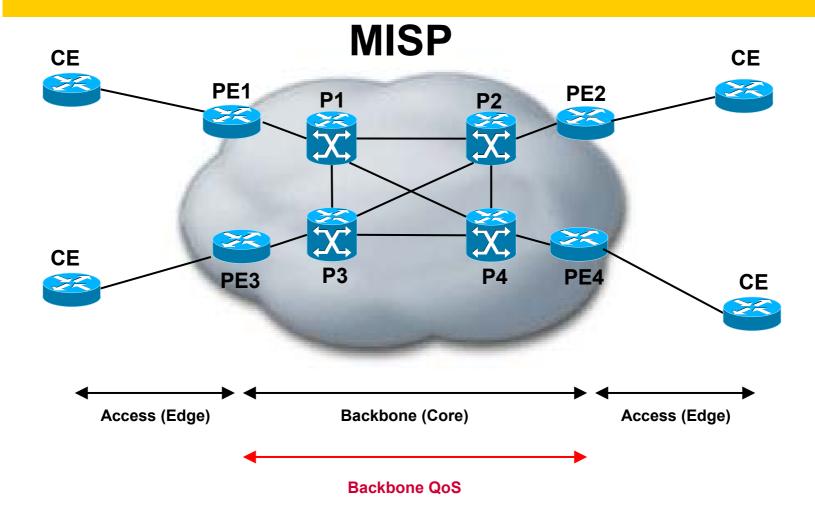
- Class 3 gets 10% of available bandwidth
- Class 2 gets 90% of available bandwidth



- IP QoS in the Backbone
  - -Between PE P / P P routers
- IP QoS in the access part
  - -Between PE CE routers



### **Backbone QoS**





#### **Backbone CoS**

- The Key is OverProvisioning Offer must be higher than Demand
- The service that traffic receives is dependent upon the ratio of traffic load to available capacity
- More Bandwidth (offer) than traffic (demand) means
  - Low loss
  - Low Latency
  - Low Jitter



### **Over-Provisioned Backbone**

### A simple rule of design:

95-Percentile (5-min average Load) <= 50% Link

which means

OverProvisioning (OP) > 2

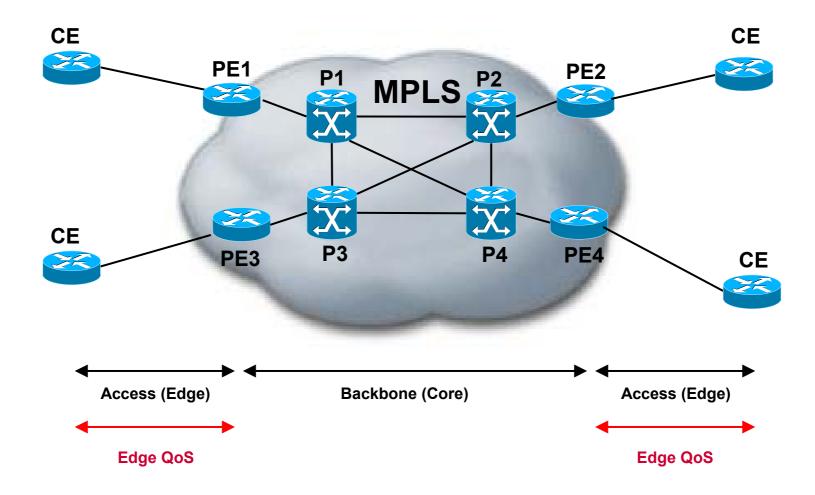


### **Backbone QoS**

- What happens if OP factor <=2 ?</li>
- Prerequisite is a constant monitoring of the backbone links.
- As soon as we approach 50% load on those links, the decision has to be made if the capacity (bandwidth) is increased (e.g. STM-1 to STM-16 / STM-16 to STM-64) or we design and deploy a MPLS aware QoS in the backbone.



# Edge (access) QoS





# Edge (access) QoS



- -Customer traffic classification
- -IP packet marking (DSCP)
- -Class based queuing (CBWFQ, CBLLQ)
- -Class-based policing for class 1
- -out-of-contract marking for class 2
- -Differentiated dropping (Class-based WRED)
- -Link fragmentation and Interleave (LFI FRF.12)

- -Class based queuing based on DSCP marking (CBWFQ, CBLLQ)
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## **Implementation**

- CoS (IP QoS) is standards (Diffserv) compliant
- CoS uses the Cisco IOS Modular QoS CLI (MQC)
- Class-based low latency queuing (CBLLQ) for Class 1
- Class-based weighted fair queuing (CBWFQ) for Class 2 and 3
- Weighted random early detection (WRED) for Class 2 and 3
- FRF.12 Link fragmentation and interleave (LFI) on links up to 768 kbps if Class 1 is present



### The invisible classes

What about management traffic (smtp, telnet, tftp etc...) and router internal traffic (keepalives, BGP etc.)?

The two invisible classes:

- MGT class (for all management traffic)
- RP class (for all internal traffic)

These classes are always configured.



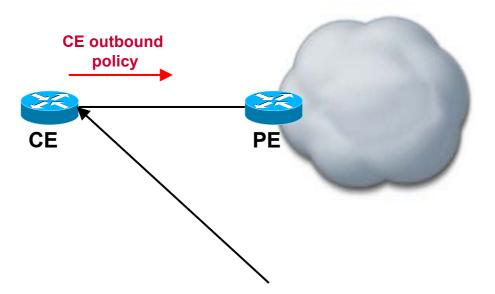
### All classes together

So in reality, we have a 5 class model.

- 3 Customer classes
  - -class 1, class 2, class 3
- 2 internal classes
  - -class MGT class RP



### **QoS Configuration on CE side**



- Decide which CoS profile (CoS classes) used on this site
- Classify customer traffic to the appropriate CoS Class

Classification is done via the MQC

=> class maps



#### Classification options

Our CoS option uses the following classification options in the class maps:

- NBAR for application
  - match protocol protocol
- access lists for IP addresses / application
  - match access-group named-access-list
- IP precedence / DSCP adapt customer marking
  - match ip precedence prec1 [prec2 [prec3 [prec4]]]
  - match ip dscp dscp1 [dscp2 [dscp3 [dscp4]]]

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#### **CE** service policy

- •What is defined in the service policy?
  - Queuing strategy
    - CBLLQ for Class 1
    - CBWFQ for the other classes
  - Intelligent dropping
    - CBWRED for Class 2 and Class 3
  - Rate limiting
    - CB policing with dropping exceeding traffic for Class 1
    - CB policing with remarking exceeding traffic for Class 2

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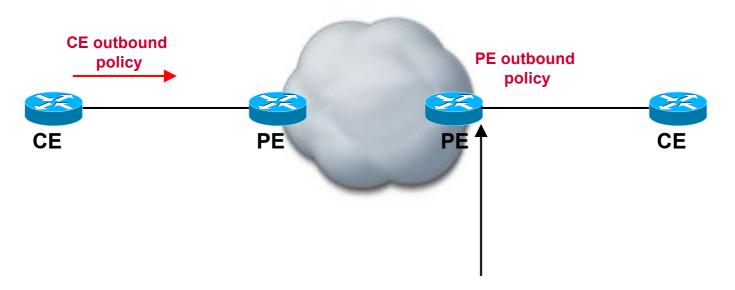
#### LFI (FRF.12) and FRTS

- In addition to the service policy, we have to use LFI (FRF.12) on links up to 768 kbps.
- •LFI with FRF.12 requires Frame Relay traffic shaping (FRTS).

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#### **QoS Configuration on PE side**



- Decide which CoS profile (CoS classes) used for this attached site
- The configuration is always the same for a certain profile and a certain link speed
  - •Packets are already correctly DSCP marked, no need to classify based on the customer specific requirements



#### **QoS Configuration on PE side**

Basically, the same mechanisms are used on the PE side as on the CE.

Due to different platforms (10K / 7600 hardware based, 7500 distributed software platform), the configuration look slightly different (nested policy-maps, no FRTS but class-based shaping, etc)

Classification is done solely based on the DSCP value of a packet



- The customer experiences problems (slow response times, time-outs, session drops) with his SQL (Oracle) and Citrix applications from time to time, especially when File transfers of his inventory systems are running.
- The traditional approach :
  - -Sell more bandwidth, although 90% of the time, bandwidth utilization is only at 50%



- Even after the bandwidth upgrade, the customer still is experiencing the same problems
  - -FTP aggressively eats up all available bandwidth

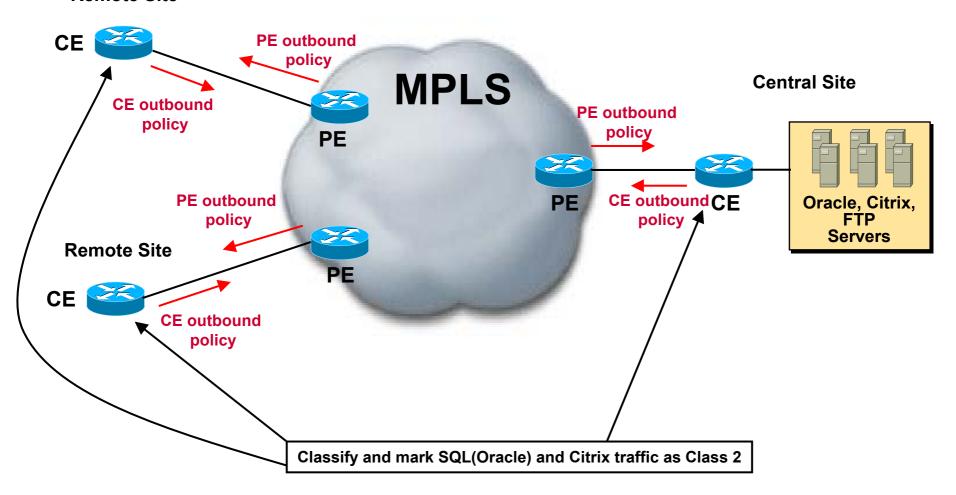
**→** The customer is very unhappy ⊗



- The sunrise CoS approach :
  - Sell customer CoS profile C with Class 2 and Class 3
  - Citrix and SQL-oracle traffic is put into Class2, the rest into Class 3

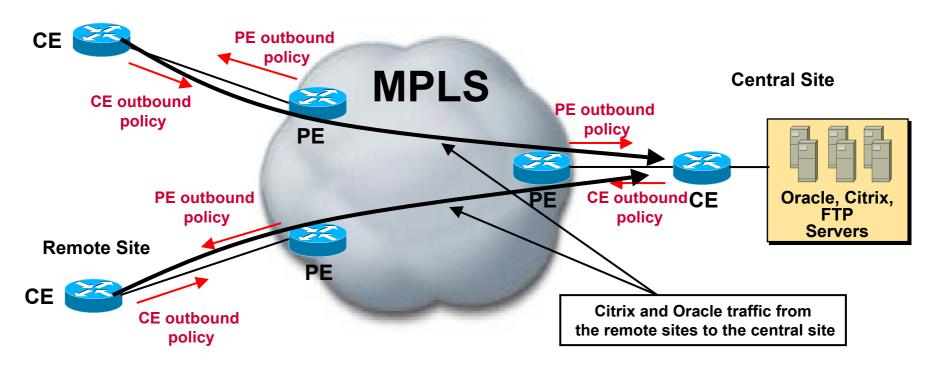








#### **Remote Site**



- •The key is appropriate PE-CE bandwidth provisioning, e.g.
  - 2 remote site à 256 kbps
  - central site 512 kbps



#### Result :

- -If there is no Oracle and Citrix traffic (e.g. during the night), the File transfer runs at full speed.
- -If there is Oracle and Citrix traffic, the File transfer traffic is throttled back to a maximum of 10% of the bandwidth, depending on the amount of Oracle and Citrix traffic.



#### Result:

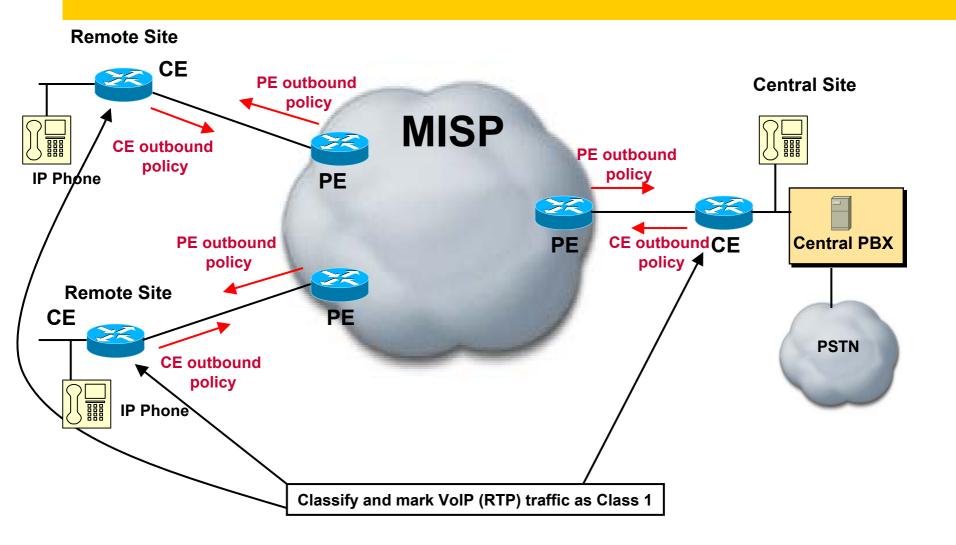
-The customer is paying less than for a bandwidth upgrade AND is getting the desired result (quality performance of Citrix and Oracle).

**→** The customer is very happy ©



- A customer has the plan to deploy IP
   Telephony (VoIP) within his VPN to get rid of
   his expensive and old PSTN infrastrucure (e.g.
   a PBX in each branch, dedicated lines etc.)
  - He is happy with the current quality and performance, just wants to add VoIP.
  - CoS is a MUST for VoIP, so he goes for profile B: Class 1 for VoIP, the rest is Class 3

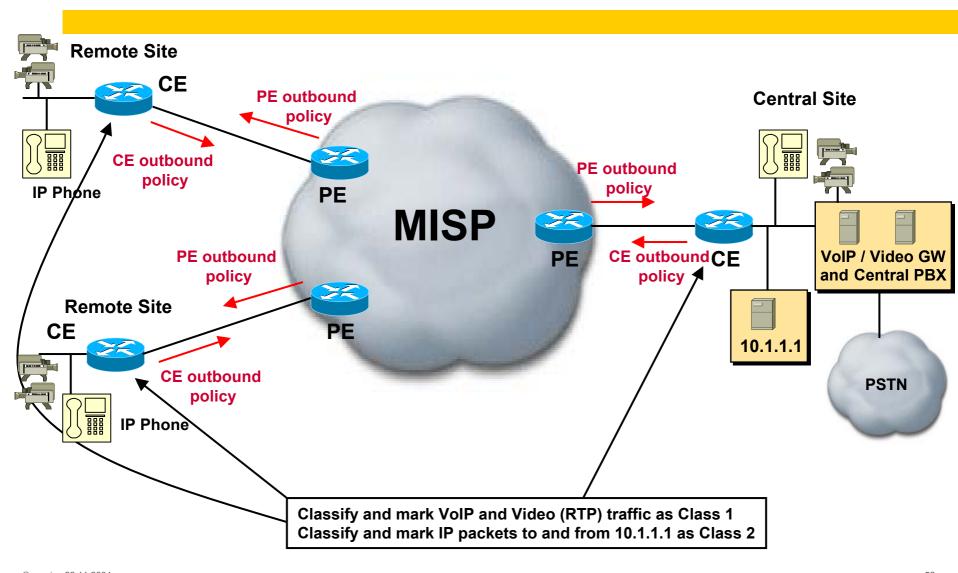






- A customer has a central server with different important applications. He wants to make sure that all of the traffic to and from this server has has priority over other data traffic.
- In addition, he plans to roll out video conferencing and/or VoIP
- The perfect customer for profile A
  - Class 1 for Voice/Video
  - Class 2 for priority data
  - Class 3 for the rest





# Class 2 "out of contract" handling



- Class 1 is fixed at 40% of the link bandwidth, excess traffic is dropped ⇒ no burst capability
  - To ensure constant low delay and jitter
- Class 2 can burst above 50% (profile A) or 90 % (profile C) if the other classes do not fully use their allocated bandwidth
- Class 3 also can burst above it's allocated bandwidth (10 % profiles A and C, 60% profile B) if room left on the link.

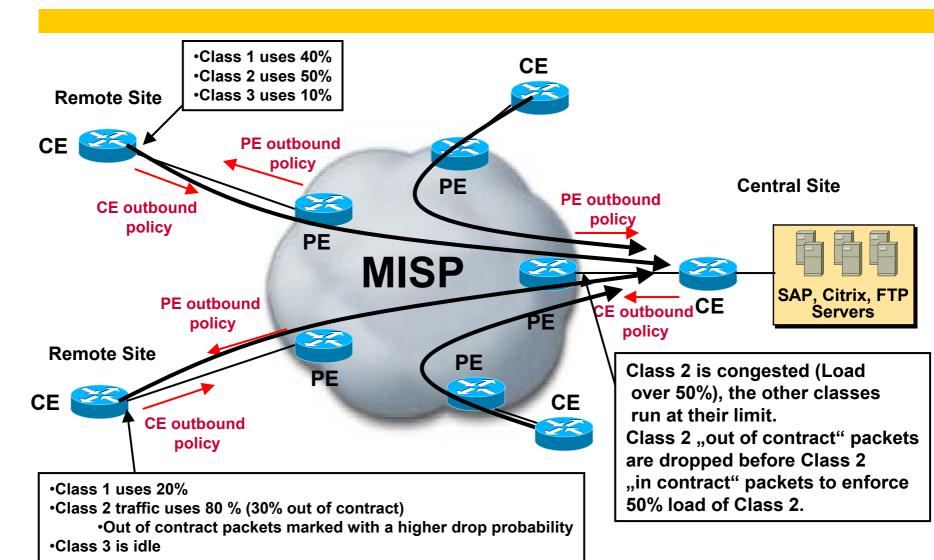
# Class 2 "out of contract" handling



- Class 3 is always treated (marked) the same way (no matter if bursting above the allocated bandwidth or "in contract")
  - Class 3 packets are always dropped first in case of congestion on a link
- Class 2 packets are treated differently if they are "in contract" or "out of contract"
  - -"Out of contract" traffic is dropped before "in contract" traffic in case of congestion

# Class 2 "out of contract" handling







#### **Common errors**

- Classification of customer traffic not working correctly
  - Wrong Class provisioning (classification of customer traffic)
- Overload of a Class (packets drops within a class)
  - Access Link bandwidth too small



#### **Common errors**

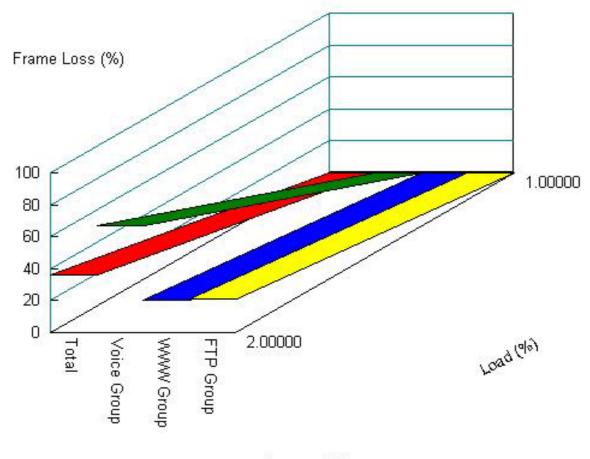
Check reporting reporting tool for Class throught and packets drops

#### alternatively

- The one and only router command which tells you all: show policy-map interface interface
- Information about
  - How many packets/bytes classified per class
  - How many packet/byte drops per class
  - and much much more…
- Remember to check CE and PE side !!!



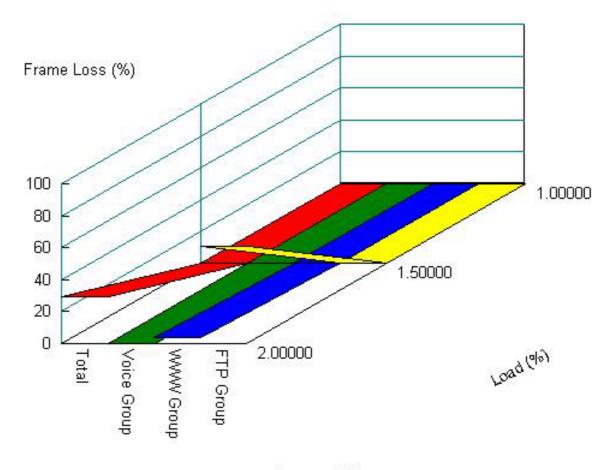
#### **Packet Loss without CoS**



Groups & Flows



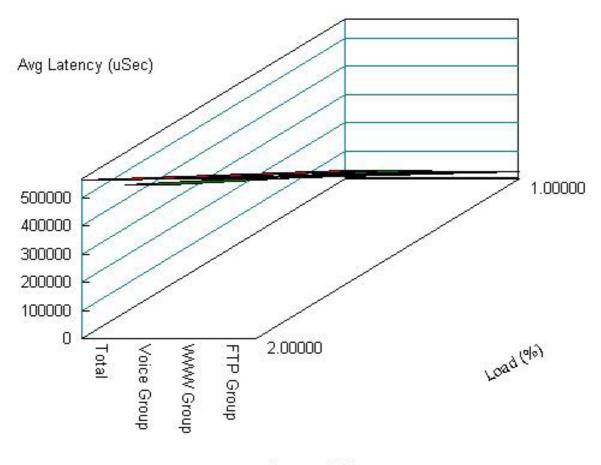
#### **Packet Loss with CoS**



Groups & Flows



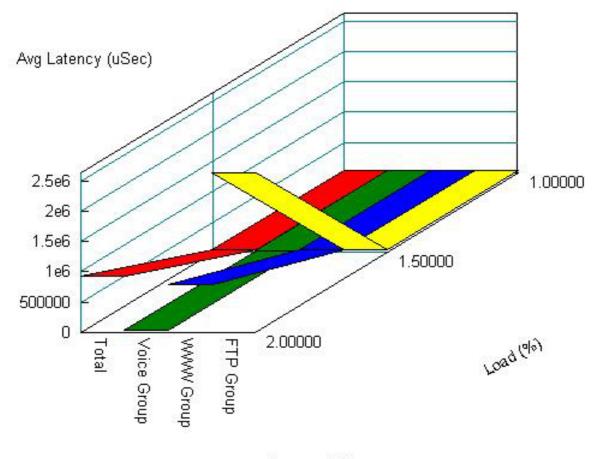
### **Delay without CoS**



Groups & Flows



### **Delay with CoS**



Groups & Flows

#### Rapid

A fatal exception 0E has occurred at 0028:C0011E36 in UXD UMM(01) + 00010E36. The current application will be terminated.

- \* Press any key to terminate the current application.
- \* Press CTRL+ALT+DEL again to restart your computer. You will lose any unsaved information in all applications.

Press any key to continue \_

